COMPUTING THE STACKELBERG/NASH EQUILIBRIA USING THE EXTRAPROXIMAL METHOD: CONVERGENCE ANALYSIS AND IMPLEMENTATION DETAILS FOR MARKOV CHAINS GAMES

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In this paper we present the extraproximal method for computing the Stackelberg/Nash equilibria in a class of ergodic controlled finite Markov chains games. We exemplify the original game formulation in terms of coupled nonlinear programming problems implementing the Lagrange principle. In addition, Tikhonov's regularization method is employed to ensure the convergence of the cost-functions to a Stackelberg/Nash equilibrium point. Then, we transform the problem into a system of equations in the proximal format. We present a two-step iterated procedure for solving the extraproximal method: (a) the first step (the extra-proximal step) consists of a "prediction" which calculates the preliminary position approximation to the equilibrium point, and (b) the second step is designed to find a "basic adjustment" of the previous prediction. The procedure is called the "extraproximal method" because of the use of an extrapolation. Each equation in this system is an optimization problem for which the necessary and efficient condition for a minimum is solved using a quadratic programming method. This solution approach provides a drastically quicker rate of convergence to the equilibrium point. We present the analysis of the convergence as well the rate of convergence of the method, which is one of the main results of this paper. Additionally, the extraproximal method is developed in terms of Markov chains for Stackelberg games. Our goal is to analyze completely a three-player Stackelberg game consisting of a leader and two followers. We provide all the details needed to implement the extraproximal method in an efficient and numerically stable way. For instance, a numerical technique is presented for computing the first step parameter (λ) of the extraproximal method. The usefulness of the approach is successfully demonstrated by a numerical example related to a pricing oligopoly model for airlines companies.

Keywords: extraproximal method, Stackelberg games, convergence analysis, Markov chains, implementation.

1. Introduction

The *extraproximal approach* (Antipin, 2005) can be considered a natural extension of the proximal and gradient optimization methods used for solving the more difficult problems of finding an equilibrium point in game theory. The simplest and most natural approach to implement the proximal method is to use a simple iteration by omitting the prediction step. However, as shown by Antipin (2005), this approach fails. A more versatile procedure would be to perform an "*ex*-

traproximal step" i.e., gathering certain information on the future development of the process. Using this information, it is possible to execute the "principle step" based on the preliminary position. It seems natural to call this procedure an "extraproximal method" (because of the use of extrapolation (Antipin, 2005)). Along with the extra-gradient technique (Poznyak, 2009), the extraproximal method can be considered a natural extension of the proximal and gradient optimization approaches to resolve more complicated game problems, such as the Stackelberg–Nash game.

Our goal is to analyze a three-player Stackelberg

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game (von Stackelberg, 1934), a leader and two followers, for a class of ergodic controllable finite Markov chains (Poznyak et al., 2000) using the extraproximal method (Antipin, 2005). The leader commits first to a strategy xin the Markov chains game. Then, the followers realize a Nash solution (Clempner and Poznyak, 2011) obtaining $\varphi_l(\cdot,\cdot|x)$ (l=1,2) as the cost functions. The optimal actions x of the leader are then chosen to minimize its payoff $\varphi_0(\cdot, \cdot | x)$. The main concern about Stackelberg games is as follows: the highest leader payoff is obtained when the followers always reply in the best possible way for the leader (this payoff is at least as high as any Nash payoff). For ground-breaking works in the Stackelberg game, see those by Bos (1986; 1991), Harris and Wiens (1980), Merril and Schneider (1966) or von Stackelberg (1934). Surveys can be found in the works of Breitmoser (2012), De Fraja and Delbono (1990), Nett (1993) or Vickers and Yarrow. (1998).

In this paper we present the extraproximal method for computing the Stackelberg/Nash equilibria in a class of ergodic controlled finite Markov chains games. We exemplify the original game formulation in terms of coupled nonlinear programming problems implementing the Lagrange principle. In addition, Tikhonov's regularization method is employed to ensure the convergence of the cost-functions to a Stackelberg/Nash equilibrium point. The Tikhonov regularization is one of the most popular approaches to solve discrete ill-posed problems represented in the form of a non-obligatory strict convex function. Then, we transform the problem to a system of equations in the proximal format.

We present a two-step iterated procedure for solving the extraproximal method: (a) the first step (the extra-proximal step) consists of a "prediction" which calculates the preliminary position approximation to the equilibrium point, and (b) the second step is designed to find a "basic adjustment" of the previous prediction. Each equation in this system is an optimization problem for which the necessary condition for a minimum is solved using a quadratic programming method, instead of the iterating projectional gradient method given by Moya and Poznyak (2009). Iterating these steps, we obtain a new quick procedure which leads to a simple and logically justified computational realization: at each iteration of both sub-steps of the procedure, the functional of the game decreases and finally converges to a Stackelberg/Nash equilibrium point (Trejo et al., 2015).

We present the analysis of the convergence as well the rate of convergence of the method. Additionally, the extraproximal method is developed in terms of Markov chains for Stackelberg games consisting of a leader and two followers. We provide all the details needed to implement the extraproximal method in an efficient and numerically stable way. For instance, a numerical method is presented for computing the first step parameter (λ) of the extraproximal method. The usefulness of the method is successfully demonstrated by a numerical example related to a pricing oligopoly model for airlines companies that choose to offer limited airline seats for specific routes.

The paper is organized as follows. The next section presents the preliminaries needed to understand the rest of the paper. The formulation of the problem for the \mathcal{N} -player game and the general format iterative version of the extraproximal method is presented in Section 2. A three-player Stackelberg game for a class of ergodic controllable finite Markov chains using the extraproximal method is given in Section 4. Section 5 includes the analysis of the converge of the extraproximal method and proves the convergence rates, which is one of the main results of this paper. A numerical example related to a pricing oligopoly model for airline companies validates the proposed extraproximal method in Section 6. Final comments are outlined in Section 7.

2. Controllable Markov chains

A controllable Markov chain (Clempner and Poznyak, 2014; Poznyak *et al.*, 2000) is a quadruple $MC = \{S, A, \Upsilon, \Pi\}$, where S is a finite set of states, $S \subset \mathbb{N}$, endowed with discrete topology; A is the set of actions, which is a metric space. For each $s \in S$, $A(s) \subset A$ is the non-empty set of admissible actions at state $s \in S$. Without loss of generality we may take $A = \bigcup_{s \in S} A(s)$; $\Upsilon = \{(s, a) | s \in S, a \in A(s)\}$ is the set of admissible state-action pairs, which is a measurable subset of $S \times A$; $\Pi(k) = [\pi_{(i,j|k)}]$ is a stationary transition controlled matrix, where

$$\pi_{(i,j|k)} \equiv P\left(s(n+1) = s_{(j)}|s(n) = s_{(i)}, a(n) = a_{(k)}\right),$$

representing the probability associated with the transition from state $s_{(i)}$ to state $s_{(j)}$ under an action $a_{(k)} \in A(s_{(i)})$ (k = 1, ..., M) at time n = 0, 1, ...

The dynamics of the game for Markov chains are described as follows. The game consists of $(\mathcal{N}+1)$ players (denoted by $l = \overline{0, \mathcal{N}}$) and begins at the initial state $s^{l}(0)$ which (as well as the states further realized by the process) is assumed to be completely measurable. Each player l is allowed to randomize, with distribution $d_{(k|i)}^{l}(n)$, over the pure action choices $a_{(k)}^{l} \in A^{l}(s_{(i)}^{l})$, $i = \overline{1, \mathcal{N}_{l}}$ and $k = \overline{1, \mathcal{M}_{l}}$. The leader corresponds to l = 0 and followers to $l = \overline{1, \mathcal{N}}$. At each fixed strategy of the leader $d_{(k|i)}^{0}(n)$, the followers make the strategy selection trying to realize a Nash equilibrium. Below we will consider only stationary strategies $d_{(k|i)}^{l}(n) = d_{(k|i)}^{l}$.

These choices induce the state distribution dynamics

$$P^{l}\left(s^{l}(n+1)=s_{(j_{l})}\right)$$
$$=\sum_{i_{l}=1}^{N_{l}}\left(\sum_{k_{l}=1}^{M_{l}}\pi^{l}_{(i_{l},j_{l}|k_{l})}d^{l}_{(k_{l}|i_{l})}\right)P^{l}\left(s^{l}(n)=s_{(i_{l})}\right)$$

In the ergodic case (when all Markov chains are ergodic for any stationary strategy $d_{(k|i)}^l$) the distributions $P^l\left(s^l(n+1)=s_{(j_l)}\right)$ exponentially quickly converge to their limits $P^l\left(s=s_{(i)}\right)$ satisfying

$$P^{l}\left(s^{l} = s_{(j_{l})}\right)$$

= $\sum_{i_{l}=1}^{N_{l}} \left(\sum_{k_{l}=1}^{M_{l}} \pi^{l}_{(i_{l},j_{l}|k_{l})} d^{l}_{(k_{l}|i_{l})}\right) P^{l}\left(s^{l} = s_{(i_{l})}\right).$ (1)

For any player l, his/her *individual rationality* is the player's cost function J^l of any fixed policy d^l , defined over all possible combinations of states and actions, and indicates the expected value when taking action a^l in state s^l and following policy d^l thereafter. The *J*-values can be expressed by

$$\mathbf{J}^{l}(c^{l}) := \sum_{i_{l},k_{l}} W^{l}_{(i_{l},k_{l})} d^{l}_{(k_{l}|i_{l})} P^{l}\left(s^{l} = s_{(i_{l})}\right), \qquad (2)$$

where

$$W_{(i_l,k_l)}^l = \sum_{j_l} J_{(i_l,j_l,k_l)}^l \pi_{(i_l,j_l|k_l)}^l,$$

the function $J_{(i_l,j_l,k_l)}^l$ being a constant at state $s_{(i_l)}$ when the action $a_{(k_l)}^l$ is applied. Then, the cost function of each player, depending on the states and actions of all participants, is given by the values $W_{(i_1,k_1;...;i_N,k_N)}^l$, so that the "average cost function" \mathbf{J}^l for each player l in the stationary regime can be expressed as

$$\mathbf{J}^{l}\left(c^{0},\ldots,c^{\mathcal{N}}\right)$$
$$:=\sum_{i_{0},k_{0}}\cdots\sum_{i_{\mathcal{N}},k_{\mathcal{N}}}W^{l}_{\left(i_{1},k_{1},\ldots,i_{\mathcal{N}},k_{\mathcal{N}}\right)}\prod_{l=0}^{\mathcal{N}}c^{l}_{\left(i_{l},k_{l}\right)},\quad(3)$$

where

$$c^{l} := \|c_{(i_{l},k_{l})}^{l}\|_{i_{l}=\overline{1,N_{l}};k_{l}=\overline{1,M_{l}}}$$

is a matrix with elements

$$c_{(i_l,k_l)}^l = d_{(k_l|i_l)}^l P^l \left(s^l = s_{(i_l)} \right)$$
(4)

satisfying

$$c^{l} \in C_{\text{adm}}^{l} = \begin{cases} c^{l} : \sum_{i_{l},k_{l}} c^{l}_{(i_{l},k_{l})} = 1, c^{l}_{(i_{l},k_{l})} \ge 0, \\ \\ \sum_{k_{l}} c^{l}_{(j_{l},k_{l})} = \sum_{i_{l},k_{l}} \pi^{l}_{(i_{l},j_{l}|k_{l})} c^{l}_{(i_{l},k_{l})}. \end{cases}$$
(5)

Notice that by (4) it follows that

$$P^{l}\left(s^{l}=s_{(i_{l})}\right) = \sum_{k_{l}} c^{l}_{(i_{l},k_{l})},$$

$$d^{l}_{(k_{l}|i_{l})} = \frac{c^{l}_{(i_{l},k_{l})}}{\sum_{k_{l}} c^{l}_{(i_{l},k_{l})}}.$$
(6)

In the ergodic case,

$$\sum_{k_l} c_{(i_l,k_l)}^l > 0$$

for all $l = \overline{0, N}$. The *individual aim* of each participant is

$$\mathbf{J}^l(c^l) \to \min_{c^l \in C^l_{\mathrm{adm}}}.$$

Obviously, here we have a conflict situation which can be resolved by the Stackelberg–Nash equilibrium concept discussed in detail below.

3. Formulation of the problem

 v^l

3.1. Stackelberg–Nash equilibrium concept. To simplify the descriptions, below let us introduce the new variables:

$$x := \operatorname{col} c^{(0)}, \quad X := C_{\operatorname{adm}}^{(0)},$$

:= col $c^{(l)}, \quad V^{l} := C_{\operatorname{adm}}^{(l)} \quad \left(l = \overline{1, \mathcal{N}}\right).$ (7)

Consider a non-zero sum game with a leader whose strategies are denoted by $x \in X$ and \mathcal{N} followers with strategies $v^l \in V^l(l = \overline{1, \mathcal{N}})$. Denote by

$$v = (v^1, \dots, v^{\mathcal{N}}) \in V := \bigotimes_{l=1}^{\mathcal{N}} V^l$$

the joint strategy of the followers. The leader is assumed to anticipate the reactions of the followers. They are trying to reach one of the Nash equilibria for any fixed strategy x of the leader, that is, to find a joint strategy $w^* = (w^{1*}, \ldots, w^{\mathcal{N}*}) \in W$ satisfying for any admissible $x \in X, w^l \in W^l$ and any $l = \overline{1, \mathcal{N}}$ the system of inequalities (the Nash condition)

$$g_l(v^l, w^{l*}|x) \le 0 \quad \text{for any } v^l \in V^l \text{ and all } l = \overline{1, \mathcal{N}},$$
$$g_l(v^l, w^{\hat{l}}|x) := \varphi_l(w^l, w^{\hat{l}}|x) - \varphi_l(v^l, w^{\hat{l}}|x),$$
(8)

where $w^{\hat{l}}$ is a strategy of the rest of the followers adjoint to v^{l} , namely,

$$w^{\hat{l}} := (v^{1}, \dots, v^{l-1}, v^{l+1}, \dots, v^{\mathcal{N}})$$
$$\in W^{\hat{l}} := \bigotimes_{m=1, \ m \neq l}^{\mathcal{N}} V^{m}$$

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$$w^l := \arg\min_{v^l \in V^l} \varphi_l\left(v^l, w^{\hat{l}} | x\right).$$

Here $\varphi_l(v^l, w^{\hat{l}}|x)$ is the cost-function of the player l, which plays the strategy $v^l \in V^l$, and the rest of the players—the strategy $w^{\hat{l}} \in W^{\hat{l}}$.

Lemma 1. The Nash equilibrium $\overline{w} \in W$ for the followers given a fixed strategy of the leader $x \in X$ (8) can be equivalently expressed in the joint format (Tanaka and Yokoyama, 1991),

$$\max_{v \in W} g(v, \bar{w}|x) \leq 0,$$

$$g(v, w|x) := \sum_{l=1}^{N} \left[\varphi_l(w^l, w^{\hat{l}}|x) - \varphi_l(v^l, w^{\hat{l}}|x) \right],$$

$$\varphi_l(w^l, w^{\hat{l}}|x) := \min_{z^l \in W^l} \varphi_l(z^l, w^{\hat{l}}|x),$$

$$v^l \in W^l, \quad w \in W := \bigotimes_{l=1}^{N} W^l.$$
(9)

Proof. Summing (8) implies (9). Conversely, taking $v^m = w^m$ for all $m \neq l$ in (9), which is valid for any admissible v^l , we obtain (8).

Notice that the condition $g(v, \bar{w}|x) \leq 0$ (9) is equivalent to

$$\max_{v \in V} \left\{ g(v, \bar{w}|x) \right\} \le 0$$

for any fixed $v \in V$ and any $x \in X$.

Let f(x,w) $(x \in X, w \in W)$ be the cost function of the leader. The functions of the followers $\varphi_l(v^l, w^{\hat{l}}|x)$ $(l = \overline{1, N})$ and the function of the leader f(x, w) are assumed to be convex in all their arguments.

Definition 1. A strategy $x^* \in X$ of the leader together with the collection $w^* \in W$ of the followers is said to be a *Stackelberg–Nash equilibrium* if

$$(x^*, w^*) \in \operatorname{Arg\,min}_{x \in X} \ \max_{v \in V, w \in W} \{f(x, w) | g(v, w | x) \le 0\}.$$
(10)

3.2. Regularized Lagrange principle application. Applying the Lagrange principle (see, for example, Poznyak *et al.*, 2000) for Definition 1, we may conclude that (10) can be rewritten as

$$(x^*, w^*) \in \operatorname{Arg\,min}_{x \in X} \max_{v \in V, w \in W, \lambda \ge 0} \mathcal{L}(x, v, w, \lambda),$$
$$\mathcal{L}(x, v, w, \lambda) := f(x, w) + \lambda g(v, w | x).$$
(11)

The approximative solution obtained by the Tikhonov regularization (see Poznyak *et al.*, 2000) is given by

$$(x_{\delta}^{*}, w_{\delta}^{*}) = \arg\min_{x \in X} \max_{v \in V, w \in W, \lambda \ge 0} \mathcal{L}_{\delta}(x, v, w, \lambda),$$
$$\mathcal{L}_{\delta}(x, v, w, \lambda) := f_{\delta}(x, w) + \lambda g_{\delta}(v, w|x) - \frac{\delta}{2}\lambda^{2},$$
(12)

where $\delta > 0$ and

$$f_{\delta}(x,w) := f(x,w) + \frac{\delta}{2} ||x||^{2},$$

$$g_{\delta}(v,w|x) = g(v,w|x) - \frac{\delta}{2} (||v||^{2} + ||w||^{2})$$
(13)

Notice that with $\delta > 0$ the functions considered become strictly convex providing the uniqueness of the discussed conditional optimization problem (12). For $\delta =$ 0 we have (11). Notice also that the Lagrange function in (12) satisfies the saddle-point condition (Poznyak, 2009), namely, for all $x \in X$ and $\lambda \ge 0$ we have

$$\mathcal{L}_{\delta}(x_{\delta}^{*}, v_{\delta}, w_{\delta}, \lambda_{\delta}) \leq \mathcal{L}_{\delta}(x_{\delta}^{*}, v_{\delta}^{*}, w_{\delta}^{*}, \lambda_{\delta}^{*}) \\ \leq \mathcal{L}_{\delta}(x_{\delta}, v_{\delta}^{*}, w_{\delta}^{*}, \lambda_{\delta}^{*}).$$

$$(14)$$

3.3. Proximal format. In the *proximal format* (see Antipin, 2005) the relation (12) can be expressed as

$$\lambda_{\delta}^{*} = \arg \max_{\lambda \ge 0} \left\{ -\frac{1}{2} \|\lambda - \lambda_{\delta}^{*}\|^{2} + \gamma \mathcal{L}_{\delta}(x_{\delta}^{*}, v_{\delta}^{*}, w_{\delta}^{*}, \lambda) \right\},$$

$$x_{\delta}^{*} = \arg \min_{x \in X} \left\{ \frac{1}{2} \|x - x_{\delta}^{*}\|^{2} + \gamma \mathcal{L}_{\delta}(x, v_{\delta}^{*}, w_{\delta}^{*}, \lambda_{\delta}^{*}) \right\},$$

$$v_{\delta}^{*} = \arg \max_{v \in V} \left\{ -\frac{1}{2} \|v - v_{\delta}^{*}\|^{2} + \gamma \mathcal{L}_{\delta}(x_{\delta}^{*}, v, w_{\delta}^{*}, \lambda_{\delta}^{*}) \right\},$$

$$w_{\delta}^{*} = \arg \max_{w \in W} \left\{ -\frac{1}{2} \|w - w_{\delta}^{*}\|^{2} + \gamma \mathcal{L}_{\delta}(x_{\delta}^{*}, v_{\delta}^{*}, w, \lambda_{\delta}^{*}) \right\},$$
(15)

where the solutions $x_{\delta}^*, v_{\delta}^*, w_{\delta}^*$ and λ_{δ}^* depend on the small parameters $\delta, \gamma > 0$.

Define the extended variables

$$\begin{split} \tilde{x} &:= \left(\begin{array}{c} x \end{array}\right) \in \tilde{X} := X, \\ \tilde{y} &:= \left(\begin{array}{c} v \\ w \\ \lambda \end{array}\right) \in \tilde{Y} := V \times W \times \mathbb{R}^+ \\ \tilde{w} &= \left(\begin{array}{c} \tilde{w}_1 \\ \tilde{w}_2 \end{array}\right) \in \tilde{X} \times \tilde{Y}, \\ \tilde{v} &= \left(\begin{array}{c} \tilde{v}_1 \\ \tilde{v}_2 \end{array}\right) \in \tilde{X} \times \tilde{Y}, \end{split}$$

and the functions

$$\hat{\mathcal{L}}_{\delta}(\tilde{x}, \tilde{y}) := \mathcal{L}_{\delta}(x, v, w, \lambda),$$

$$\Psi_{\delta}(\tilde{w}, \tilde{v}) := \tilde{\mathcal{L}}_{\delta}(\tilde{w}_1, \tilde{v}_2) - \tilde{\mathcal{L}}_{\delta}(\tilde{v}_1, \tilde{w}_2).$$

For $\tilde{w}_1 = \tilde{x}$, $\tilde{w}_2 = \tilde{y}$, $\tilde{v}_1 = \tilde{v}_1^* = \tilde{x}_\delta^*$ and $\tilde{v}_2 = \tilde{v}_2^* = \tilde{y}_\delta^*$, we have

$$\Psi_{\delta}(\tilde{w}, \tilde{v}) := \hat{\mathcal{L}}_{\delta}(\tilde{x}, \tilde{y}_{\delta}^*) - \hat{\mathcal{L}}_{\delta}(\tilde{x}_{\delta}^*, \tilde{y}).$$

In these variables the relation (15) can be represented in a *"short format"* as

$$\tilde{v}^* = \arg\min_{\tilde{w}\in\tilde{X}\times\tilde{Y}} \left\{ \frac{1}{2} \|\tilde{w} - \tilde{v}^*\|^2 + \gamma \Psi_{\delta}(\tilde{w}, \tilde{v}^*) \right\}.$$
(16)

3.4. Extraproximal method. The *extraproximal method* for the conditional optimization problems (12) was suggested by Antipin (2005). We design the method for the static Stackelberg–Nash game in a general format.

The general format iterative version (n = 0, 1, ...) of the extraproximal method with some fixed admissible initial values $(x_0 \in X, v_0 \in V, w_0 \in W, \text{ and } \lambda_0 \ge 0)$ is as follows:

1. The first half-step (prediction):

$$\overline{\lambda}_{n} = \arg\min_{\lambda \geq 0} \left\{ \frac{1}{2} \|\lambda - \lambda_{n}\|^{2} - \gamma \mathcal{L}_{\delta}(x_{n}, v_{n}, w_{n}, \lambda) \right\},$$

$$\overline{x}_{n} = \arg\min_{x \in X} \left\{ \frac{1}{2} \|x - x_{n}\|^{2} + \gamma \mathcal{L}_{\delta}(x, v_{n}, w_{n}, \bar{\lambda}_{n}) \right\},$$

$$\overline{v}_{n} = \arg\min_{v \in V} \left\{ \frac{1}{2} \|v - v_{n}\|^{2} - \gamma \mathcal{L}_{\delta}(x_{n}, v, w_{n}, \bar{\lambda}_{n}) \right\},$$

$$\overline{w}_{n} = \arg\min_{w \in W} \left\{ \frac{1}{2} \|w - w_{n}\|^{2} - \gamma \mathcal{L}_{\delta}(x_{n}, v_{n}, w, \bar{\lambda}_{n}) \right\},$$
(17)

or, in the extended variables,

$$\hat{v}_n = \arg\min_{\tilde{w}\in \tilde{X}\times \tilde{Y}} \left\{ \frac{1}{2} \|\tilde{w} - \tilde{v}_n\|^2 + \gamma \Psi_{\delta}(\tilde{w}, \tilde{v}_n) \right\}.$$
(18)

2. The second (basic) half-step,

$$\lambda_{n+1} = \arg\min_{\lambda\geq 0} \left\{ \frac{1}{2} \|\lambda - \lambda_n\|^2 - \gamma \mathcal{L}_{\delta}(\bar{x}_n, \bar{v}_n, \bar{w}_n, \lambda) \right\},$$

$$x_{n+1} = \arg\min_{x\in X} \left\{ \frac{1}{2} \|x - x_n\|^2 + \gamma \mathcal{L}_{\delta}(x, \bar{v}_n, \bar{w}_n, \bar{\lambda}_n) \right\},$$

$$v_{n+1} = \arg\min_{v\in V} \left\{ \frac{1}{2} \|v - v_n\|^2 - \gamma \mathcal{L}_{\delta}(\bar{x}_n, v, \bar{w}_n, \bar{\lambda}_n) \right\},$$

$$w_{n+1} = \arg\min_{w\in W} \left\{ \frac{1}{2} \|w - w_n\|^2 - \gamma \mathcal{L}_{\delta}(\bar{x}_n, \bar{v}_n, w, \bar{\lambda}_n) \right\},$$
(19)

or, in the "short format",

$$\tilde{v}_{n+1} = \arg\min_{\tilde{w}\in\tilde{X}\times\tilde{Y}} \left\{ \frac{1}{2} \|\tilde{w} - \tilde{v}_n\|^2 + \gamma \Psi_{\delta}(\tilde{w}, \hat{v}_n) \right\}.$$
(20)

Remark 1. The direct one-step procedure, when $\lambda_{n+1} = \overline{\lambda}_n$, does not work (see the counter-example of Antipin (2005)).

4. Markov format for the extraproximal method

We consider a three-player Stackelberg game $(l = \overline{1, N})$. Player l should be understood to refer to players in a general expression of the game. In the remainder of this paper, the leader is designated as player (0) and the followers as players (1) and (2). Numbers are used to refer to players in a specific expression. We also assume that the number of strategies and actions that this description can take is finite and fixed. Here we will apply the iterative quadratic method providing a significantly quicker rate of convergence.

4.1. Cost functions and notation. The individual cost-functions of the followers and the leader are defined as follows.

Cost-function for the leader:

$$\mathbf{J}^{(0)}(c^{(1)}, c^{(2)}, c^{(0)}) = \sum_{i_1, k_1} \sum_{i_2, k_2} \sum_{i_0, k_0} W^{(0)}_{(i_1, k_1; i_2, k_2; i_0, k_0)} c^{(1)}_{(i_1, k_1)} c^{(2)}_{(i_2, k_2)} c^{(0)}_{(i_0, k_0)}.$$
(21)

Cost-functions for the followers:

Follower 1: $\mathbf{J}^{(1)}(c^{(1)}, c^{(2)}, c^{(0)})$

$$=\sum_{i_1,k_1}\sum_{i_2,k_2}\sum_{i_0,k_0}W^{(1)}_{(i_1,k_1;i_2,k_2;i_0,k_0)}c^{(1)}_{(i_1,k_1)}c^{(2)}_{(i_2,k_2)}c^{(0)}_{(i_0,k_0)}.$$
(22)

Follower 2:

$$\mathbf{J}^{(2)}(c^{(1)}, c^{(2)}, c^{(0)}) = \sum_{i_1, k_1} \sum_{i_2, k_2} \sum_{i_0, k_0} W^{(2)}_{(i_1, k_1; i_2, k_2; i_0, k_0)} c^{(1)}_{(i_1, k_1)} c^{(2)}_{(i_2, k_2)} c^{(0)}_{(i_0, k_0)}.$$
(23)

For $n = 0, 1, \ldots$, let us define the vectors

$$x_n = \operatorname{col} c^{(0)}(n),$$
$$v_n = \begin{pmatrix} \operatorname{col} c^{(1)}(n) \\ \operatorname{col} c^{(2)}(n) \end{pmatrix},$$

$$w_n = w^{\hat{l}} = \left(\begin{array}{c} \operatorname{col} c^{(\hat{1})}(n) \\ \operatorname{col} c^{(\hat{2})}(n) \end{array}\right),$$

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$$\begin{split} w^{l} &= \left(\begin{array}{c} \operatorname{col} \mathring{c}^{(1)}\left(n\right)\\ \operatorname{col} \mathring{c}^{(2)}\left(n\right) \end{array}\right) \\ &= \left(\begin{array}{c} \operatorname{col} \arg\min_{z \in C_{adm}^{(1)}} \mathbf{J}^{(1)}\left(z, c^{(\hat{2})}\left(n\right), x_{n}\right)\\ \operatorname{col} \arg\min_{z \in C_{adm}^{(2)}} \mathbf{J}^{(2)}\left(c^{(\hat{1})}\left(n\right), z, x_{n}\right) \end{array}\right). \end{split}$$

Let us introduce for simplicity the following notation:

$$W_{(i_1,k_1;i_2,k_2;i_0,k_0)}^{(l)} = W^{(l)}$$

$$c_{(i_1,k_l)}^{(l)} = c^{(l)}$$

$$c^{(\hat{l})}(i_l,k_l) = c^{(\hat{l})}$$

$$\hat{c}_{(i_l,k_l)}^{(l)} = \hat{c}^{(l)}$$

$$\overline{c}_{(i_l,k_l)}^{(l)} = \overline{c}^{(l)}$$

For the leader we have

$$f(x,w) = \sum_{i_1,k_1} \sum_{i_2,k_2} \sum_{i_0,k_0} W^{(0)} \mathring{c}^{(1)} \mathring{c}^{(2)} c^{(0)}.$$
 (24)

From

$$\begin{split} \mathbf{J}^{(1)}\left(\mathring{c}^{(1)}, c^{(\hat{2})} | c^{(0)}\right) &\leq \mathbf{J}^{(1)}\left(c^{(1)}, c^{(\hat{2})} | c^{(0)}\right), \\ \mathbf{J}^{(2)}\left(\mathring{c}^{(2)}, c^{(\hat{1})} | c^{(0)}\right) &\leq \mathbf{J}^{(2)}\left(c^{(2)}, c^{(\hat{1})} | c^{(0)}\right), \end{split}$$

we have

$$\varphi_{l}\left(w^{l}, w^{\hat{l}}|x\right)$$

$$= \mathbf{J}^{(1)}\left(\mathring{c}^{(1)}, c^{(\hat{2})}|c^{(0)}\right) + \mathbf{J}^{(2)}\left(\mathring{c}^{(2)}, c^{(\hat{1})}|c^{(0)}\right)$$

$$= \sum_{i_{1},k_{1}} \sum_{i_{2},k_{2}} \sum_{i_{0},k_{0}} \left[W^{(1)}\mathring{c}^{(1)}c^{(\hat{2})}c^{(0)} + W^{(2)}c^{(\hat{1})}\mathring{c}^{(2)}c^{(0)}\right],$$
(25)

as well as

$$\varphi_{l}\left(v^{l}, w^{\hat{l}}|x\right) = \mathbf{J}^{(1)}\left(c^{(1)}, c^{(\hat{2})}|c^{(0)}\right) + \mathbf{J}^{(2)}\left(c^{(2)}, c^{(\hat{1})}|c^{(0)}\right) = \sum_{i_{1}, k_{1}} \sum_{i_{2}, k_{2}} \sum_{i_{0}, k_{0}} \left[W^{(1)}c^{(1)}c^{(\hat{2})}c^{(0)} + W^{(2)}c^{(\hat{1})}c^{(2)}c^{(0)}\right].$$
(26)

For computing the Nash equilibrium point, we define the function g(v, w, x) as follows:

$$g(v, w|x) = \sum_{i_1, k_1} \sum_{i_2, k_2} \sum_{i_0, k_0} \left[W^{(1)} \mathring{c}^{(1)} c^{(\hat{2})} c^{(0)} + W^{(2)} c^{(\hat{1})} \mathring{c}^{(2)} c^{(0)} - W^{(1)} c^{(1)} c^{(\hat{2})} c^{(0)} - W^{(2)} c^{(\hat{1})} c^{(2)} c^{(0)} \right].$$
(27)

4.2. Tikhonov's regularization method. The regularized problem for (24) is defined as follows:

$$f_{\delta}(x,w) = \sum_{i_1,k_1} \sum_{i_2,k_2} \sum_{i_0,k_0} W^{(0)} \mathring{c}^{(1)} \mathring{c}^{(2)} c^{(0)} + \frac{\delta}{2} \|c^{(0)}\|^2,$$

and for (27) we have

$$g_{\delta}(v, w|x) = \sum_{i_{1},k_{1}} \sum_{i_{2},k_{2}} \sum_{i_{0},k_{0}} \left[W^{(1)} \mathring{c}^{(1)} c^{(\hat{2})} c^{(0)} + W^{(2)} c^{(\hat{1})} \mathring{c}^{(2)} c^{(0)} - W^{(1)} c^{(1)} c^{(\hat{2})} c^{(0)} - W^{(2)} c^{(\hat{1})} c^{(2)} c^{(0)} \right] \\ - \frac{\delta}{2} \left(\left\| \begin{array}{c} c^{(1)} \\ c^{(2)} \end{array} \right\|^{2} + \left\| \begin{array}{c} c^{(\hat{1})} \\ c^{(\hat{2})} \end{array} \right\|^{2} \right).$$
(28)

4.3. Lagrange principle. Let us apply the Lagrange principle (Poznyak, 2009) with the Lagrange function

$$\begin{split} \mathfrak{L}_{\delta}(x,v,w,\lambda) &= f_{\delta}(x,w) + \lambda g_{\delta}(v,w,x) - \frac{\delta}{2}\lambda^{2} \\ &= \sum_{i_{1},k_{1}} \sum_{i_{2},k_{2}} \sum_{i_{0},k_{0}} \left[W^{(0)} \mathring{c}^{(1)} \mathring{c}^{(2)} c^{(0)} \\ &+ \lambda \left(W^{(1)} \mathring{c}^{(1)} c^{(\hat{2})} c^{(0)} + W^{(2)} c^{(\hat{1})} \mathring{c}^{(2)} c^{(0)} \\ &- W^{(1)} c^{(1)} c^{(\hat{2})} c^{(0)} - W^{(2)} c^{(\hat{1})} c^{(2)} c^{(0)} \right) \right] \\ &- \lambda \frac{\delta}{2} \left(\left\| \begin{array}{c} c^{(1)} \\ c^{(2)} \end{array} \right\|^{2} + \left\| \begin{array}{c} c^{(\hat{1})} \\ c^{(\hat{2})} \end{array} \right\|^{2} \right) + \frac{\delta}{2} \| c^{(0)} \|^{2} - \frac{\delta}{2} \lambda^{2}. \end{split}$$

4.4. Extrapoximal method for Markov chains. The procedure for solving the Stackelberg/Nash equilibrium point for the extraproximal method consists of the following "iterative rules" implementation.

1. First half-step: 1.a. For

$$\overline{\lambda}_n = \arg\min_{\lambda \ge 0} \left\{ \frac{1}{2} (\lambda - \lambda_n)^2 - \gamma \mathfrak{L}_{\delta}(x_n, v_n, w_n, \lambda) \right\},\$$

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we have

$$\begin{split} \overline{\lambda}_n &= \left[\lambda_n + \gamma \sum_{i_1, k_1} \sum_{i_2, k_2} \sum_{i_0, k_0} \left(W^{(1)} \mathring{c}^{(1)} c^{(2)}(n) c^{(0)}(n) \right. \\ &+ W^{(2)} c^{(1)}(n) \mathring{c}^{(2)} c^{(0)}(n) \\ &- W^{(1)} c^{(1)}(n) c^{(2)}(n) c^{(0)}(n) \\ &- W^{(2)} c^{(1)}(n) c^{(2)}(n) c^{(0)}(n) \right) \\ &- \gamma \frac{\delta}{2} \left(\left\| \begin{array}{c} c^{(1)}(n) \\ c^{(2)}(n) \end{array} \right\|^2 + \left\| \begin{array}{c} c^{(1)}(n) \\ c^{(2)}(n) \end{array} \right\|^2 \right) \right]^+ \\ &\times \frac{1}{1 + \gamma \delta}. \end{split}$$

1.b. For

$$\overline{x}_n = \arg\min_{x \in X} \left\{ \frac{1}{2} \|x - x_n\|^2 + \gamma \mathfrak{L}_{\delta}(x, v_n, w_n, \overline{\lambda}_n) \right\},\$$

we have

$$\begin{split} \overline{x}_n &= \arg\min_{x\in X} \left\{ \left(\frac{1}{2} + \gamma \frac{\delta}{2}\right) \|c^{(0)}\|^2 \\ &+ \gamma \sum_{i_1,k_1} \sum_{i_2,k_2} \sum_{i_0,k_0} \left[W^{(0)} \dot{c}^{(1)} \dot{c}^{(2)} c^{(0)} \\ &+ \overline{\lambda}_n \left(W^{(1)} \dot{c}^{(1)} c^{(\hat{2})}(n) c^{(0)} \\ &+ W^{(2)} c^{(\hat{1})}(n) \dot{c}^{(2)} c^{(0)} - W^{(1)} c^{(1)}(n) c^{(\hat{2})}(n) c^{(0)} \\ &- W^{(2)} c^{(\hat{1})}(n) c^{(2)}(n) c^{(0)} \right) \right] \\ &- c^{(0)T}(n) c^{(0)} + \frac{1}{2} \|c^{(0)}(n)\|^2 \\ &- \gamma \overline{\lambda}_n \frac{\delta}{2} \left(\left\| \begin{array}{c} c^{(1)}(n) \\ c^{(2)}(n) \end{array} \right\|^2 + \left\| \begin{array}{c} c^{(\hat{1})}(n) \\ c^{(\hat{2})}(n) \end{array} \right\|^2 \right) \\ &- \gamma \frac{\delta}{2} \overline{\lambda}_n^2 \right\}. \end{split}$$

1.c. For

$$\overline{v}_n = \arg\min_{v\in V} \left\{ \frac{1}{2} \|v - v_n\|^2 - \gamma \mathfrak{L}_{\delta}(x_n, v, w_n, \overline{\lambda}_n) \right\},\$$

we have

$$\overline{v}_n = \arg\min_{v \in V} \left\{ \left(\frac{1}{2} + \gamma \overline{\lambda}_n \frac{\delta}{2} \right) \left\| \begin{array}{c} c^{(1)} \\ c^{(2)} \end{array} \right\|^2 \right\}$$

$$- \left(\begin{array}{c} c^{(1)} \\ c^{(2)} \end{array} \right)^{T} \left(\begin{array}{c} c^{(1)} \\ c^{(2)} \\ n \end{array} \right)$$

$$+ \gamma \overline{\lambda}_{n} \sum_{i_{1},k_{1}} \sum_{i_{2},k_{2}} \sum_{i_{0},k_{0}} \left[W^{(1)} c^{(1)} c^{(\hat{2})} (n) c^{(0)} (n) \right.$$

$$+ W^{(2)} c^{(\hat{1})} (n) c^{(2)} c^{(0)} (n) - W^{(1)} \mathring{c}^{(1)} c^{(\hat{2})} (n) c^{(0)} (n)$$

$$- W^{(2)} c^{(\hat{1})} (n) \mathring{c}^{(2)} c^{(0)} (n) \right]$$

$$- \gamma \sum_{i_{1},k_{1}} \sum_{i_{2},k_{2}} \sum_{i_{0},k_{0}} W^{(0)} \mathring{c}^{(1)} \mathring{c}^{(2)} c^{(0)} (n)$$

$$- \gamma \frac{\delta}{2} \| c^{(0)} (n) \|^{2} + \frac{1}{2} \left\| \begin{array}{c} c^{(1)} (n) \\ c^{(2)} (n) \end{array} \right\|^{2}$$

$$+ \gamma \overline{\lambda}_{n} \frac{\delta}{2} \left\| \begin{array}{c} c^{(\hat{1})} (n) \\ c^{(\hat{2})} (n) \end{array} \right\|^{2} + \gamma \frac{\delta}{2} \overline{\lambda}_{n}^{2} \right\}.$$

1.d. For

we have

$$\begin{split} \overline{w}_{n} &= \arg\min_{w\in W} \left\{ \left(\frac{1}{2} + \gamma \overline{\lambda}_{n} \frac{\delta}{2}\right) \left\| \begin{array}{c} c^{(\hat{1})} \\ c^{(\hat{2})} \end{array} \right\|^{2} \\ &- \left(\begin{array}{c} c^{(\hat{1})} \\ c^{(\hat{2})} \end{array} \right)^{T} \left(\begin{array}{c} c^{(\hat{1})}(n) \\ c^{(\hat{2})}(n) \end{array} \right) \\ &- \gamma \overline{\lambda}_{n} \sum_{i_{1},k_{1}} \sum_{i_{2},k_{2}} \sum_{i_{0},k_{0}} \left[W^{(1)} \mathring{c}^{(1)} c^{(\hat{2})} c^{(0)}(n) \\ &+ W^{(2)} c^{(\hat{1})} \mathring{c}^{(2)} c^{(0)}(n) - W^{(1)} c^{(1)}(n) c^{(\hat{2})} c^{(0)}(n) \\ &- W^{(2)} c^{(\hat{1})} c^{(2)}(n) c^{(0)}(n) \right] \\ &- \gamma \sum_{i_{1},k_{1}} \sum_{i_{2},k_{2}} \sum_{i_{0},k_{0}} W^{(0)} \mathring{c}^{(1)} \mathring{c}^{(2)} c^{(0)}(n) \\ &- \gamma \frac{\delta}{2} \| c^{(0)}(n) \|^{2} + \frac{1}{2} \left\| \begin{array}{c} c^{(\hat{1})}(n) \\ c^{(\hat{2})}(n) \end{array} \right\|^{2} \\ &+ \gamma \overline{\lambda}_{n} \frac{\delta}{2} \left\| \begin{array}{c} c^{(1)}(n) \\ c^{(2)}(n) \end{array} \right\|^{2} + \gamma \frac{\delta}{2} \overline{\lambda}_{n}^{2} \right\}. \end{split}$$

 $\overline{w}_n = \arg\min_{w \in W} \left\{ \frac{1}{2} \|w - w_n\|^2 - \gamma \mathfrak{L}_{\delta}(x_n, v_n, w, \overline{\lambda}_n) \right\},\$

2. Second half-step:

2.a. For

$$\lambda_{n+1} = \arg\min_{\lambda \ge 0} \bigg\{ \frac{1}{2} (\lambda - \lambda_n)^2 - \gamma \mathfrak{L}_{\delta}(\overline{x}_n, \overline{w}_n, \overline{w}_n, \lambda) \bigg\},\,$$

we have

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$$\begin{aligned} x_{n+1} &= \arg\min_{x \in X} \left\{ \frac{1}{2} \|x - x_n\|^2 + \gamma \mathfrak{L}_{\delta}(x, \overline{v}_n, \overline{w}_n, \overline{\lambda}_n) \right. \\ \text{we obtain} \\ &= \arg\min_{x \in X} \left\{ \left(\frac{1}{2} + \gamma \frac{\delta}{2} \right) \|c^{(0)}\|^2 \\ &+ \gamma \sum_{i_1, k_1} \sum_{i_2, k_2} \sum_{i_0, k_0} \left[W^{(0)} \mathring{c}^{(1)} \mathring{c}^{(2)} c^{(0)} \\ &+ \overline{\lambda}_n \left(W^{(1)} \mathring{c}^{(1)} \overline{c}^{(2)}(n) c^{(0)} \right) \right] \end{aligned}$$

$$\times \frac{1}{1 + \gamma \delta}.$$
2.b. For
$$x_{n+1} = \arg\min_{x \in X} \left\{ \frac{1}{2} \|x - x_n\|^2 + \gamma \mathfrak{L}_{\delta}(x, \overline{v}_n, \overline{w}_n, \overline{\lambda}_n) \right\},$$
we obtain
$$x_{n+1}$$

 $-\gamma \frac{\delta}{2} \left(\left\| \frac{\overline{c}^{(1)}(n)}{\overline{c}^{(2)}(n)} \right\|^2 + \left\| \frac{\overline{c}^{(\hat{1})}(n)}{\overline{c}^{(\hat{2})}(n)} \right\|^2 \right) \right]^+$

 $\lambda_{n+1} = \left| \lambda_n + \gamma \sum_{i_1, k_2} \sum_{i_2, k_3} \sum_{i_3, k_3} \left(W^{(1)} \mathring{c}^{(1)} \overline{c}^{(2)}(n) \overline{c}^{(0)}(n) \right) \right|^2$

 $+ W^{(2)}\overline{c}^{(\hat{1})}(n)\mathring{c}^{(2)}\overline{c}^{(0)}(n)$ $-W^{(1)}\overline{c}^{(1)}(n)\overline{c}^{(2)}(n)\overline{c}^{(0)}(n)$ $-W^{(2)}\overline{c}^{(\hat{1})}(n)\overline{c}^{(2)}(n)\overline{c}^{(0)}(n)\Big)$

$$= \arg\min_{x \in X} \left\{ \left(\frac{1}{2} + \gamma \frac{0}{2} \right) \| c^{(0)} \|^{2} + \gamma \sum_{i_{1},k_{1}} \sum_{i_{2},k_{2}} \sum_{i_{0},k_{0}} \left[W^{(0)} \mathring{c}^{(1)} \mathring{c}^{(2)} c^{(0)} + \overline{\lambda}_{n} \left(W^{(1)} \mathring{c}^{(1)} \overline{c}^{(2)} (n) c^{(0)} + W^{(2)} \overline{c}^{(1)} (n) \mathring{c}^{(2)} (n) c^{(0)} - W^{(1)} \overline{c}^{(1)} (n) \overline{c}^{(2)} (n) c^{(0)} - W^{(2)} \overline{c}^{(1)} (n) \overline{c}^{(2)} (n) c^{(0)} \right) \right] \\ - c^{(0)T} (n) c^{(0)} + \frac{1}{2} \| c^{(0)} (n) \|^{2} \\ - \gamma \overline{\lambda}_{n} \frac{\delta}{2} \left(\left\| \frac{\overline{c}^{(1)} (n)}{\overline{c}^{(2)} (n)} \right\|^{2} + \left\| \frac{\overline{c}^{(1)} (n)}{\overline{c}^{(2)} (n)} \right\|^{2} \right) \\ - \gamma \frac{\delta}{2} \overline{\lambda}_{n}^{2} \right\}.$$

2.c. For

$$v_{n+1} = \arg\min_{v\in V} \left\{ \frac{1}{2} \|v - v_n\|^2 - \gamma \mathfrak{L}_{\delta}(\overline{x}_n, v, \overline{w}_n, \overline{\lambda}_n) \right\},\$$

we get

$$v_{n+1} = \arg\min_{v \in V} \left\{ \left(\frac{1}{2} + \gamma \overline{\lambda}_n \frac{\delta}{2} \right) \left\| \begin{array}{c} c^{(1)} \\ c^{(2)} \end{array} \right\|^2 \right\}$$

we have

$$w_{n+1} = \arg\min_{w \in W} \left\{ \frac{1}{2} \|w - w_n\|^2 - \gamma \mathfrak{L}_{\delta}(\overline{x}_n, \overline{v}_n, w, \overline{\lambda}_n) \right\},$$

 $w_{n+1} = \arg\min_{w \in W} \left\{ \left(\frac{1}{2} + \gamma \overline{\lambda}_n \frac{\delta}{2} \right) \left\| \begin{array}{c} c^{(\hat{1})} \\ c^{(\hat{2})} \end{array} \right\|^2 \right\}$

 $- \left(\begin{array}{c} c^{(\hat{1})} \\ c^{(\hat{2})} \end{array}\right)^T \left(\begin{array}{c} c^{(\hat{1})}(n) \\ c^{(\hat{2})}(n) \end{array}\right)$

 $+ W^{(2)} c^{(\hat{1})}_{(i_1,k_1)} \mathring{c}^{(2)} \overline{c}^{(0)}(n)$ $-W^{(1)}\overline{c}^{(1)}(n)c^{(\hat{2})}\overline{c}^{(0)}(n)$

 $- W^{(2)} c^{(\hat{1})} \overline{c}^{(2)}(n) \overline{c}^{(0)}(n) \Big]$

 $-\gamma \overline{\lambda}_n \sum_{i_1,k_1} \sum_{i_2,k_2} \sum_{i_0,k_0} \left[W^{(1)} \mathring{c}^{(1)} c^{(\hat{2})} \overline{c}^{(0)}(n) \right]$

 $-\gamma \sum_{i_1,k_1} \sum_{i_2,k_2} \sum_{i_0,k_0} W^{(0)} \mathring{c}^{(1)} \mathring{c}^{(2)} \overline{c}^{(0)}(n)$

 $-\gamma \frac{\delta}{2} \|\overline{c}^{(0)}(n)\|^2 + \frac{1}{2} \left\| \begin{array}{c} c^{(\hat{1})}(n) \\ c^{(\hat{2})}(n) \end{array} \right\|^2$

 $+\gamma\overline{\lambda}_n\frac{\delta}{2} \left\| \frac{\overline{c}^{(1)}(n)}{\overline{c}^{(2)}(n)} \right\|^2 +\gamma\frac{\delta}{2}\overline{\lambda}_n^2 \right\}.$

4.5. Quadratic programming solver. Let us restrict the expression (3) for the admissible set and the stationary

case with the following constraint.

$$w_{n+1} = \arg\min_{w \in W} \left\{ \frac{1}{2} \|w - w_n\|^2 - \gamma \mathfrak{L}_{\delta}(\overline{x}_n, \overline{v}_n, w, \overline{\lambda}_n) \right\},\$$

2.d. For

$$- \left(\begin{array}{c} c^{(2)} \end{array} \right) \left(\begin{array}{c} c^{(2)} (n) \end{array} \right) \\ + \gamma \overline{\lambda}_{n} \sum_{i_{1},k_{1}} \sum_{i_{2},k_{2}} \sum_{i_{0},k_{0}} \left[W^{(1)} c^{(1)} \overline{c}^{(\hat{2})}(n) \overline{c}^{(0)}(n) \\ + W^{(2)} \overline{c}^{(\hat{1})}(n) c^{(2)} \overline{c}^{(0)}(n) - W^{(1)} \dot{c}^{(1)} \overline{c}^{(\hat{2})}(n) \overline{c}^{(0)}(n) \\ - W^{(2)} \overline{c}^{(\hat{1})}(n) \dot{c}^{(2)} \overline{c}^{(0)}(n) \right] \\ - \gamma \sum_{i_{1},k_{1}} \sum_{i_{2},k_{2}} \sum_{i_{0},k_{0}} W^{(0)} \dot{c}^{(1)} \dot{c}^{(2)} \overline{c}^{(0)}(n) - \gamma \frac{\delta}{2} \| \overline{c}^{(0)}(n) \|^{2} \\ + \frac{1}{2} \left\| \begin{array}{c} c^{(1)}(n) \\ c^{(2)}(n) \end{array} \right\|^{2} + \gamma \overline{\lambda}_{n} \frac{\delta}{2} \left\| \begin{array}{c} \overline{c}^{(\hat{1})}(n) \\ \overline{c}^{(\hat{2})}(n) \end{array} \right\|^{2} + \gamma \frac{\delta}{2} \overline{\lambda}_{n}^{2} \right\}.$$

 $\begin{pmatrix} c^{(1)} \end{pmatrix}^T \begin{pmatrix} c^{(1)}(n) \end{pmatrix}$

Each matrix
$$d_{(i_l|k_l)}^{(l)}$$
 represents a stationary mixed-strategy that belongs to the simplex
$$\int d_{(i_l|k_l)}^{(l)} \in \mathbb{R}^{N_l M_l} \text{ for } d_{(i_l|k_l)}^{(l)} \ge 0,$$

$$S^{(N_l M_l)} := \begin{cases} d_{(i_l|k_l)}^{(l)} \in \mathbb{R}^{N_l M_l} \text{ for } d_{(i_l|k_l)}^{(l)} \ge \\ \\ \text{where } \sum_{k_l=1}^{M_l} d_{(i_l|k_l)}^{(l)} = 1. \end{cases}$$

A necessary and sufficient condition for $d_{(i_l|k_l)}^{(l)}$ to be a Stackelberg/Nash equilibrium point is the solution of the following QPP (quadratic programming problem):

$$\mathbf{J}^{l}\left(d^{0}_{(i_{0},k_{0})},d^{1}_{(i_{1},k_{1})},\ldots,d^{\mathcal{N}}_{(i_{\mathcal{N}},k_{\mathcal{N}})}\right) \to \min_{d^{(l)}_{(i_{l}|k_{l})}} \quad (29)$$

subject to

$$d_{(i_l|k_l)}^{(l)} \in S^{(N_l M_l)}.$$

We consider the problem

$$\frac{1}{2}X^{\top}HX + C^{\top}X \to \min_X \,, \tag{30}$$

 $AX \le b, \quad AX = b,$

where $0 \leq X \leq 1$, $A^l \in \mathbb{R}^{N_l + 1 \times (N_l M_l)}$ and $b^l \in \mathbb{R}^{N_l + 1}$.

The vector $C^T = C_{(N|M)}$ is defined as

$$C_{(1|1)} = W_{(1|1)}^{l},$$

$$C_{(2|1)} = W_{(2|1)}^{l},$$

$$\vdots$$

$$C_{(N|1)} = W_{(N|1)}^{l},$$

$$\vdots$$

$$C_{(1|M_{l})} = W_{(1|M_{l})}^{l},$$

$$\vdots$$

$$C_{(N_{l}|M_{l})} = W_{(N_{l}|M_{l})}^{l}$$

that is, $C_{(N|M)}$ is equal to $W^l_{(i_1,k_1,\ldots,i_{\mathcal{N}},k_{\mathcal{N}})}$ ordered by columns.

The vector b^l is defined as

$$b^l = (0, \dots, 0, 1)^T \in \mathbb{R}^{N_l + 1}.$$

We will construct the matrix $A^l \in \mathbb{R}^{N_l \times (N_l M_l)}$ using the ergodicity constraints defined in (5),

$$0 = \sum_{k_l=1}^{M_l} \left(\sum_{i_l=1}^{N_l} \pi_{(i_l j_l | k_l)} c_{(i_l | k_l)}^{(l)} - c_{(j_l | k_l)}^{(l)} \right).$$
(31)

Then, we have

$$j_{l} = 1 \sum_{k_{l}=1}^{M_{l}} \left(\sum_{i_{l}=1}^{N_{l}} \pi_{(i_{l}j_{l}|k_{l})} c_{(i_{l}|k_{l})}^{(l)} - c_{(1|k_{l})}^{(l)} \right) = 0,$$

$$j_{l} = 2 \sum_{k_{l}=1}^{M_{l}} \left(\sum_{i_{l}=1}^{N_{l}} \pi_{(i_{l}j_{l}|k_{l})} c_{(i_{l}|k_{l})}^{(l)} - c_{(2|k_{l})}^{(l)} \right) = 0,$$

:

$$j_l = N_l \sum_{k_l=1}^{M_l} \left(\sum_{i_l=1}^{N_l} \pi_{(i_l j_l | k_l)} c_{(i_l | k_l)}^{(l)} - c_{(N_l | k_l)}^{(l)} \right) = 0.$$

Developing the formulas of (31) for \mathcal{N} players and multiplying by $c_{(i_1|k_1)}^{(l)}$ for each component (31), we have

$$BL^{(l)} = \left[\pi^{(l)}_{(j_l, i_l|k_l)} - \delta_{(j_l, i_l)} \right]_{j_l = \overline{1, N_l}}_{i_l = \overline{1, N_l}} \\ = \left[\begin{array}{ccc} \pi^{(l)}_{(1, 1|1)} - 1 & \dots & \pi^{(l)}_{(N_l, 1|M_l)} \\ \dots & \dots & \dots \\ \pi^{(l)}_{(N_l, 1|1)} & \dots & \pi^{(l)}_{(N_l, N_l|M_l)} - 1 \end{array} \right],$$

where $\delta_{(j_l,i_l)}$ is Kronecker's delta. Then $A_{\text{eq}} \in \mathbb{R}(\sum_{l=1}^{N} N_l + (\mathcal{N})) \times (\sum_{l=1}^{N} N_l \times M_l)$ is defined as follows:

$$A_{\rm eq} = \begin{bmatrix} BL^{(1)} & 0 & \dots & 0 \\ 0 & BL^{(2)} & \dots & 0 \\ \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & BL^{(M_l)} \\ [1] & 0 & \dots & 0 \\ 0 & [1] & \dots & 0 \\ 0 & 0 & \dots & 0 \\ 0 & 0 & 0 & [1] \end{bmatrix},$$

where

$$[1] = (1, \dots, 1)^{\mathsf{T}} \in \mathbb{R}^{(N_l)}$$

 $b_{\mathrm{eq}} \in \mathbb{R}^{\left(\sum_{l=1}^{\mathcal{N}} N_l + (\mathcal{N})\right)}$

and

is

$$b_{\mathrm{eq}} = \left(\begin{array}{cccccccc} 0 & \dots & 0 & 1 & \dots & 1 \end{array} \right)^{\top}.$$

5. Convergence analysis

5.1. Auxiliary results. Let us prove the following auxiliary results.

Lemma 2. Let f(z) be a convex function defined on a convex set Z. If z^* is a minimizer of function

$$\varphi(z) = \frac{1}{2} ||z - x||^2 + \alpha f(z)$$
 (32)

on Z with fixed x, then f(z) satisfies the inequality

$$\frac{1}{2} \|z^* - x\|^2 + \alpha f(z^*)$$

$$\leq \frac{1}{2} \|z - x\|^2 + \alpha f(z) - \frac{1}{2} \|z - z^*\|^2. \quad (33)$$

Proof. By the necessary condition for a minimum at z^* ,

$$\langle z^* - x + \alpha \nabla f(z^*), z - z^* \rangle \ge 0,$$

and the convexity property of f(z), expressed as

$$f(z) \ge f(z^*) + \langle \nabla f(z^*), z - z^* \rangle,$$

we derive

$$0 \le \langle z^* - x + \alpha \nabla f(z^*), z - z^* \rangle$$

= $\langle z^* - x, z - z^* \rangle + \alpha [f(z) - f(z^*)]$.

Using this inequality in the identity

$$\begin{aligned} \frac{1}{2} \|z - x\|^2 \\ &= \frac{1}{2} \|z - z^*\|^2 + \langle z^* - x, z - z^* \rangle + \frac{1}{2} \|z^* - x\|^2, \end{aligned}$$

we get (33). The lemma is proven.

Lemma 3. If all partial derivatives of $\tilde{\mathcal{L}}_{\delta}(\tilde{x}, \tilde{y})$ satisfy the Lipschitz condition with positive constant C_0 , then the following Lipschitz-type condition holds:

$$\begin{aligned} \| [\Psi_{\delta}(\tilde{w}+h,\tilde{v}+g) - \Psi_{\delta}(\tilde{w},\tilde{v}+g)] \\ - [\Psi_{\delta}(\tilde{w}+h,\tilde{v}) - \Psi_{\delta}(\tilde{w},\tilde{v})] \| \le C_0 \|h\| \|g\| \quad (34) \end{aligned}$$

valid for any $\tilde{w}, h, \tilde{v}, g \in \tilde{X} \times \tilde{Y}$.

Proof. By Lagrange's formula

$$f(x+h) - f(x) = \int_0^1 \langle \nabla f(x+th), h \rangle \,\mathrm{d}t$$

we have

$$\begin{split} \| [\Psi_{\delta}(\tilde{w}+h,\tilde{v}+g) - \Psi_{\delta}(\tilde{w},\tilde{v}+g)] \\ &- [\Psi_{\delta}(\tilde{w}+h,\tilde{v}) - \Psi_{\delta}(\tilde{w},\tilde{v})] \| \\ = \| \int_{0}^{1} \langle \nabla \Psi_{\delta}(\tilde{w}+th,\tilde{v}+g) - \nabla \Psi_{\delta}(\tilde{w}+th,\tilde{v}),h \rangle \, \mathrm{d}t \| \\ &\leq \int_{0}^{1} | \langle \nabla \Psi_{\delta}(\tilde{w}+th,\tilde{v}+g) - \nabla \Psi_{\delta}(\tilde{w}+th,\tilde{v}),h \rangle | \, \mathrm{d}t \\ &\leq \int_{0}^{1} C_{0} \|h\| \|g\| \, \mathrm{d}t \leq C_{0} \|h\| \|k\|, \end{split}$$

which proves the lemma.

5.2. Main convergence theorem. The following theorem presents the convergence conditions of (17)–(19) and gives an estimate of its rate of convergence.

Theorem 1. Assume that the problem (16) has a solution. Let $\tilde{\mathcal{L}}_{\delta}(\tilde{x}, \tilde{y})$ be differentiable in \tilde{x} and \tilde{y} , whose partial derivative with respect to \tilde{y} satisfies the Lipschitz condition with positive constant C. Then, for any $\delta \in (0, 1)$, there exists a small enough

$$\gamma_0 = \gamma_0(\delta) < C := \min\left\{\frac{1}{\sqrt{2}C_0}, \frac{1 + \sqrt{1 + 2C_0^2}}{2C_0^2}\right\}$$

such that, for any $0 < \gamma \leq \gamma_0$, the sequence $\{\tilde{v}_n\}$, generated by the equivalent extraproximal procedure (18)–(20), monotonically converges in norm with a geometric progression rate $q \in (0, 1)$ to one of the equilibrium points \tilde{v}^* , i.e.,

$$\|\widetilde{v}_n - \widetilde{v}^*\|^2 \le q^n \|\widetilde{v}_0 - \widetilde{v}^*\|^2, \tag{35}$$

where

$$q = 1 + \frac{4(\delta\gamma)^2}{1 + 2\delta\gamma - 2\gamma^2 C^2} - 2\delta\gamma.$$
(36)

Proof. Taking in (33) $\alpha = \gamma$ and

$$\begin{split} z &= \tilde{w}, \quad x = \tilde{v}_n, \quad z^* = \hat{v}_n, \\ f(z) &= \Psi_{\delta}(\tilde{w}, \tilde{v}_n), \quad f(z^*) = \Psi_{\delta}(\hat{v}_n, \tilde{v}_n) \end{split}$$

we obtain

$$\begin{split} & \frac{1}{2} \| \hat{v}_n - \widetilde{v}_n \|^2 + \gamma \Psi_{\delta}(\hat{v}_n, \tilde{v}_n) \\ & \leq \frac{1}{2} \| \tilde{w} - \tilde{v}_n \|^2 + \gamma \Psi_{\delta}(\tilde{w}, \tilde{v}_n) - \frac{1}{2} \| \tilde{w} - \hat{v}_n \|^2. \end{split}$$
(37)

Again setting in (33) $\alpha = \gamma$ and

$$\begin{aligned} z &= \tilde{w}, \quad x = \tilde{v}_n, \quad z^* = \tilde{v}_{n+1}, \\ f(z) &= \Psi_{\delta}(\tilde{w}, \hat{v}_n), \quad f(z^*) = \Psi_{\delta}(\tilde{v}_{n+1}, \hat{v}_n) \end{aligned}$$

we get

$$\begin{split} & \frac{1}{2} \| \tilde{v}_{n+1} - \tilde{v}_n \|^2 + \gamma \Psi_{\delta}(\tilde{v}_{n+1}, \hat{v}_n) \\ & \leq \frac{1}{2} \| \tilde{w} - \tilde{v}_n \|^2 + \gamma \Psi_{\delta}(\tilde{w}, \hat{v}_n) - \frac{1}{2} \| \tilde{w} - \tilde{v}_{n+1} \|^2. \end{split}$$
(38)

Selecting $\tilde{w} = \tilde{v}_{n+1}$ in (37) and $\tilde{w} = \hat{v}_n$ in (38), we obtain

$$\begin{aligned} \frac{1}{2} \| \hat{v}_n - \tilde{v}_n \|^2 + \gamma \Psi_{\delta}(\hat{v}_n, \tilde{v}_n) \\ &\leq \frac{1}{2} \| \tilde{v}_{n+1} - \tilde{v}_n \|^2 + \gamma \Psi_{\delta}(\tilde{v}_{n+1}, \tilde{v}_n) \\ &- \frac{1}{2} \| \tilde{v}_{n+1} - \hat{v}_n \|^2, \end{aligned}$$
(39)
$$\begin{aligned} \frac{1}{2} \| \tilde{v}_{n+1} - \tilde{v}_n \|^2 + \gamma \Psi_{\delta}(\tilde{v}_{n+1}, \hat{v}_n) \\ &\leq \frac{1}{2} \| \hat{v}_n - \tilde{v}_n \|^2 + \gamma \Psi_{\delta}(\hat{v}_n, \hat{v}_n) - \frac{1}{2} \| \hat{v}_n - \tilde{v}_{n+1} \|^2. \end{aligned}$$
(40)

Adding (39) with (40) and using (34) for

$$\begin{split} \tilde{w}+h &= \tilde{v}_{n+1}, \quad \tilde{w} = \hat{v}_n, \quad \tilde{v}+g = \tilde{v}_n, \quad \tilde{v} = \hat{v}_n, \\ h &= \tilde{v}_{n+1} - \hat{v}_n, \quad g = \tilde{v}_n - \hat{v}_n, \end{split}$$

we finally conclude that

$$\begin{split} \|\tilde{v}_{n+1} - \hat{v}_n\|^2 &\leq \gamma [\Psi_{\delta}(\tilde{v}_{n+1}, \tilde{v}_n) - \Psi_{\delta}(\hat{v}_n, \tilde{v}_n)] \\ &- \gamma [\Psi_{\delta}(\tilde{v}_{n+1}, \hat{v}_n) - \Psi_{\delta}(\hat{v}_n, \hat{v}_n)] \\ &\leq \gamma C \|\tilde{v}_{n+1} - \hat{v}_n\| \|\tilde{v}_n - \hat{v}_n\|, \end{split}$$

which implies

$$\|\tilde{v}_{n+1} - \hat{v}_n\| \le \gamma C \|\tilde{v}_n - \hat{v}_n\|.$$
(41)

Now, taking $\tilde{w} = \tilde{v}_{n+1}$ in (37) and $\tilde{w} = \tilde{v}_{\delta}^*$ in (38), we get

$$\begin{split} &\frac{1}{2} \| \hat{v}_n - \tilde{v}_n \|^2 + \gamma \Psi_{\delta}(\hat{v}_n, \tilde{v}_n) \\ &\leq \frac{1}{2} \| \tilde{v}_{n+1} - \tilde{v}_n \|^2 + \gamma \Psi_{\delta}(\tilde{v}_{n+1}, \tilde{v}_n) \\ &\quad - \frac{1}{2} \| \tilde{v}_{n+1} - \hat{v}_n \|^2, \end{split}$$

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$$\begin{split} &\frac{1}{2} \|\tilde{v}_{n+1} - \tilde{v}_n\|^2 + \gamma \Psi_{\delta}(\tilde{v}_{n+1}, \hat{v}_n) \\ &\leq \frac{1}{2} \|\tilde{v}_{\delta}^* - \tilde{v}_n\|^2 + \gamma \Psi_{\delta}(\tilde{v}_{\delta}^*, \hat{v}_n) - \frac{1}{2} \|\tilde{v}_{\delta}^* - \tilde{v}_{n+1}\|^2. \end{split}$$

Adding these two inequalities and multiplying by two yields

$$\begin{split} \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n+1}\|^{2} + \|\tilde{v}_{n+1} - \hat{v}_{n}\|^{2} + \|\hat{v}_{n} - \tilde{v}_{n}\|^{2} \\ - 2\gamma \Psi_{\delta}(\tilde{v}_{\delta}^{*}, \hat{v}_{n}) + 2\gamma [\Psi_{\delta}(\tilde{v}_{n+1}, \hat{v}_{n}) \\ + \Psi_{\delta}(\hat{v}_{n}, \tilde{v}_{n}) - \Psi_{\delta}(\tilde{v}_{n+1}, \tilde{v}_{n})] \leq \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n}\|^{2} \end{split}$$

Adding and subtracting the term $\Psi_{\delta}(\hat{v}_n, \hat{v}_n)$, we have

$$\begin{split} \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n+1}\|^{2} + \|\tilde{v}_{n+1} - \hat{v}_{n}\|^{2} + \|\hat{v}_{n} - \tilde{v}_{n}\|^{2} \\ &+ 2\gamma \left[\Psi_{\delta}(\hat{v}_{n}, \hat{v}_{n}) - \Psi_{\delta}(\tilde{v}_{\delta}^{*}, \hat{v}_{n})\right] + 2\gamma \left[\Psi_{\delta}(\tilde{v}_{n+1}, \hat{v}_{n}) \right] \\ &- \Psi_{\delta}(\hat{v}_{n}, \hat{v}_{n}) + \Psi_{\delta}(\hat{v}_{n}, \tilde{v}_{n}) - \Psi_{\delta}(\tilde{v}_{n+1}, \tilde{v}_{n})\right] \\ &\leq \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n}\|^{2}. \end{split}$$

Using (34) with $\tilde{w} + h = \tilde{v}_{n+1}$, $\tilde{w} = \hat{v}_n$, $\tilde{v} + k = \tilde{v}_n$ and $\tilde{v} = \hat{v}_n$, we have $h = \tilde{v}_{n+1} - \hat{v}_n$ and $k = \tilde{v}_n - \hat{v}_n$, and the inequality above becomes

$$\begin{split} \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n+1}\|^{2} + \|\tilde{v}_{n+1} - \hat{v}_{n}\|^{2} + \|\hat{v}_{n} - \tilde{v}_{n}\|^{2} \\ + 2\gamma \left[\Psi_{\delta}(\hat{v}_{n}, \hat{v}_{n}) - \Psi_{\delta}(\tilde{v}_{\delta}^{*}, \hat{v}_{n})\right] \\ - 2\gamma C \|\tilde{v}_{n+1} - \hat{v}_{n}\| \|\tilde{v}_{n} - \hat{v}_{n}\| \leq \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n}\|^{2} \end{split}$$

Applying (41) to the last term on the left-hand side and in view of the strict convexity property of Ψ_{δ} given by

$$\Psi_{\delta}(\hat{v}_n, \hat{v}_n) - \Psi_{\delta}(\tilde{v}^*_{\delta}, \hat{v}_n) \ge \delta \|\hat{v}_n - \tilde{v}^*_{\delta}\|^2,$$

we get

$$\begin{split} \|\tilde{v}_{\delta}^{*}-\tilde{v}_{n+1}\|^{2}+\|\tilde{v}_{n+1}-\hat{v}_{n}\|^{2}+2\gamma\delta\|\hat{v}_{n}-\tilde{v}_{\delta}^{*}\|^{2} \\ + \left(1-2\gamma^{2}C^{2}\right)\|\tilde{v}_{n}-\hat{v}_{n}\|^{2} \leq \|\tilde{v}_{\delta}^{*}-\tilde{v}_{n}\|^{2}. \end{split}$$

Applying the identity $2\langle a - c, c - b \rangle = ||a - b||^2 - ||a - c||^2 - ||c - b||^2$ with $a = \hat{v}_n$, $b = \tilde{v}_{\delta}^*$ and $c = \tilde{v}_n$ to the left-hand side of the last inequality, we have

$$\begin{split} \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n+1}\|^{2} + \|\tilde{v}_{n+1} - \hat{v}_{n}\|^{2} + \left(1 - 2\gamma^{2}C^{2}\right)\|\tilde{v}_{n} - \hat{v}_{n}\|^{2} \\ + 2\gamma\delta[2\langle\hat{v}_{n} - \tilde{v}_{n}, \tilde{v}_{n} - \tilde{v}_{\delta}^{*}\rangle + \|\tilde{v}_{n} - \hat{v}_{n}\|^{2} + \|\tilde{v}_{n} - \tilde{v}_{\delta}^{*}\|^{2} \\ &= \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n+1}\|^{2} + \|\tilde{v}_{n+1} - \hat{v}_{n}\|^{2} \\ + \left(1 + 2\gamma\delta - 2\gamma^{2}C^{2}\right)\|\tilde{v}_{n} - \hat{v}_{n}\|^{2} \\ + 4\gamma\delta\langle\hat{v}_{n} - \tilde{v}_{n}, \tilde{v}_{n} - \tilde{v}_{\delta}^{*}\rangle \\ + 2\gamma\delta\|\tilde{v}_{n} - \tilde{v}_{\delta}^{*}\|^{2} \leq \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n}\|^{2} \,. \end{split}$$

Defining $d = 1 + 2\gamma\delta - 2\gamma^2C^2$ and completing the square form of the third and fourth terms yields

$$\begin{split} \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n+1}\|^{2} + \|\tilde{v}_{n+1} - \hat{v}_{n}\|^{2} + d\|\tilde{v}_{n} - \hat{v}_{n}\|^{2} \\ + 4\gamma\delta\langle\hat{v}_{n} - \tilde{v}_{n}, \tilde{v}_{n} - \tilde{v}_{\delta}^{*}\rangle + \frac{(2\gamma\delta)^{2}}{d}\|\tilde{v}_{n} - \tilde{v}_{\delta}^{*}\|^{2} \\ - \frac{(2\gamma\delta)^{2}}{d}\|\tilde{v}_{n} - \tilde{v}_{\delta}^{*}\|^{2} + 2\gamma\delta\|\tilde{v}_{n} - \tilde{v}_{\delta}^{*}\|^{2} \le \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n}\|^{2} \end{split}$$

and

$$\begin{split} \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n+1}\|^{2} + \|\tilde{v}_{n+1} - \hat{v}_{n}\|^{2} \\ + \|\sqrt{d}(\tilde{v}_{n} - \hat{v}_{n}) + \frac{2\gamma\delta}{\sqrt{d}}(\tilde{v}_{n} - \tilde{v}_{\delta}^{*})\|^{2} \\ \leq \left(1 - 2\gamma\delta + \frac{(2\gamma\delta)^{2}}{d}\right) \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n}\|^{2}, \end{split}$$

finally implying that

$$\begin{aligned} \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n+1}\|^{2} &\leq q \|\tilde{v}_{\delta}^{*} - \tilde{v}_{n}\|^{2} \\ &\leq q^{n+1} \|\tilde{v}_{\delta}^{*} - \tilde{v}_{0}\|^{2} \xrightarrow[n \to \infty]{} \end{aligned}$$

with

$$q = 1 - 2\gamma\delta + \frac{(2\gamma\delta)^2}{d} \in (0,1)$$

The theorem is proven.

6. Numerical example

We consider a pricing oligopoly model (Karpowicz, 2012; Nowak and Romaniuk, 2013) with a leader and two followers. For instance, let us consider AA, one of the most profitable airlines globally, which has always had the reputation of a leader and industry challenger. Let us also consider two local flying companies called AI and AII as the followers of AA on the market. AI and AII were forced to immediately start competing with international airlines for routes, getting access to airports, securing flight slots and landing rights, and attracting a new customer base.

We suppose that each consumer buys at most one ticket. It is a common practice for airlines to price discriminate, i.e., consumers coming to the market at different times have valuations inversely related to the length of time between the purchase and the flight. This is the case of business people who decide to travel at the last moment, and pay much more money for getting a ticket. On the other hand, people who plan to travel a year ahead pay significantly less money. The linear demand assumes a continuum of people between these extremes.

The dynamics of the game are as follows: (a) AA is the first mover and manages all its advantage attracting the highest value customers who pay a uniformly high price, (b) AI and AII price discriminates over the residual demand. The firms can choose to offer limited airline seats for a specific route: high, medium, low, or none at all. The market price of the seats decreases with increasing total quantity offered by the companies. If the companies decide to offer a high quantity of airline seats, the price collapses so that profits drop to zero. Figure 1 describes the states and actions corresponding to the Markov chain of the pricing problem.

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Fig. 1. Pricing Markov chain.

Following Eqn. (2), the individual utilities are as follows:

high 0 22 25 0	medium 30 39 37 0	low 43 47 43 0	none 69 71 61 0
high	medium	low	none
0	22	24	0
30	39	37	0
43	47	43	0
59	41	55	0
high 0 28 37 54	medium 24 39 45 39	low 27 39 43 59	none 3 5 10 0
	0 22 25 0 high 0 30 43 59 high 0 28	$\begin{array}{cccc} & 30 \\ 22 & 39 \\ 25 & 37 \\ 0 & 0 \\ \end{array}$ high medium $0 & 22 \\ 30 & 39 \\ 43 & 47 \\ 59 & 41 \\ \end{array}$ high medium $0 & 24 \\ 28 & 39 \\ 37 & 45 \\ \end{array}$	$\begin{array}{cccccccccccccccccccccccccccccccccccc$

The transition matrices are

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$\pi^{AI}_{(i,j 1)} =$	$\begin{bmatrix} 0.3288 \\ 0.1308 \\ 0.1394 \\ 0.5846 \end{bmatrix}$	$\begin{array}{c} 0.2152 \\ 0.3639 \\ 0.1225 \\ 0.1538 \end{array}$	$\begin{array}{c} 0.3739 \\ 0.4474 \\ 0.4843 \\ 0.0937 \end{array}$	$\begin{array}{c} 0.0821 \\ 0.0578 \\ 0.2538 \\ 0.1679 \end{array}$,
$\pi^{AI}_{(i,j 2)} =$	$\begin{bmatrix} 0.4525 \\ 0.0755 \\ 0.0373 \\ 0.2750 \end{bmatrix}$	$\begin{array}{c} 0.3901 \\ 0.2378 \\ 0.1512 \\ 0.1673 \end{array}$	$\begin{array}{c} 0.1061 \\ 0.3888 \\ 0.7654 \\ 0.3104 \end{array}$	$\begin{array}{c} 0.0513 \\ 0.2979 \\ 0.0461 \\ 0.2474 \end{array}$,

$\pi^{AII}_{(i,j 1)} =$	$\begin{bmatrix} 0.2564 \\ 0.0115 \\ 0.2662 \\ 0.0513 \end{bmatrix}$	$\begin{array}{c} 0.7202 \\ 0.3802 \\ 0.0333 \\ 0.5530 \end{array}$	$\begin{array}{c} 0.0005 \\ 0.2628 \\ 0.3895 \\ 0.3145 \end{array}$	$\begin{array}{c} 0.0230\\ 0.3455\\ 0.3111\\ 0.0812 \end{array}$,
$\pi^{AII}_{(i,j 2)} =$	$\begin{bmatrix} 0.0077 \\ 0.0758 \\ 0.2946 \\ 0.0611 \end{bmatrix}$	$\begin{array}{c} 0.0217 \\ 0.3909 \\ 0.1857 \\ 0.2346 \end{array}$	$\begin{array}{c} 0.0346 \\ 0.3933 \\ 0.1726 \\ 0.5648 \end{array}$	$\begin{array}{c} 0.9360 \\ 0.1401 \\ 0.3470 \\ 0.1395 \end{array}$,
$\pi^{AA}_{(i,j 1)} =$	$\begin{bmatrix} 0.0239\\ 0.2961\\ 0.0975\\ 0.2028 \end{bmatrix}$	$\begin{array}{c} 0.8486 \\ 0.3267 \\ 0.1078 \\ 0.2719 \end{array}$	$\begin{array}{c} 0.0471 \\ 0.2228 \\ 0.6152 \\ 0.3546 \end{array}$	$\begin{array}{c} 0.0804 \\ 0.1543 \\ 0.1794 \\ 0.1707 \end{array}$,
$\pi^{AA}_{(i,j 2)} =$	$\begin{bmatrix} 0.3579 \\ 0.5572 \\ 0.1751 \\ 0.0829 \end{bmatrix}$	0.0507 0.3785 0.2076 0.2980	$\begin{array}{c} 0.2801 \\ 0.0120 \\ 0.4572 \\ 0.3511 \end{array}$	$\begin{array}{c} 0.3113 \\ 0.0523 \\ 0.1602 \\ 0.2681 \end{array}$	

Each company knows how the increased offer lowers airline seats price and their profits. Strategic dependence is a common practice in the global airline industry. In this case, AA does this so it can predict the strategic pricing movements of other airlines. By anticipating the predicted response of the followers, AA is able to make a strategy to become and stay successful. Given $\delta = 14.3$, $\gamma =$

0.008 and applying the extraproximal method for Markov chains (see Section 4.4), we obtain $c_{(i|k)}$ for the leader (see Fig. 3) and $c_{(i|k)}$ for the followers (see Figs. 4 and 5). Figure 2 shows the convergence of the parameter λ . Then, applying (6), we have that the strategies needed to converge to a Stackelberg equilibrium are as follows:

$$d^{AA} = \begin{bmatrix} 0.8995 & 0.1005\\ 0.0656 & 0.9344\\ 0.2172 & 0.7828\\ 0.2827 & 0.7173 \end{bmatrix},$$
 (42)

$$d^{AI} = \begin{bmatrix} 0.1077 & 0.8923\\ 0.1311 & 0.8689\\ 0.9035 & 0.0965\\ 0.8349 & 0.1651 \end{bmatrix},$$
(43)

$$d^{AII} = \begin{bmatrix} 0.7331 & 0.2669\\ 0.9214 & 0.0786\\ 0.1230 & 0.8770\\ 0.8856 & 0.1144 \end{bmatrix}.$$
 (44)

Air travel has become a commodity and major routes are saturated with fierce competition. Low-cost carriers have significantly influenced consumer behavior for cheap price bargains among regular travelers and increasingly among business travelers. AA has consistently been one of the most profitable airlines globally, and has always had the reputation of a leader and industry challenger fixing its strategies (42) only on the most lucrative routes. AI, looking at the strategies of AA, decided on a fully branded product/service differentiation strategy (43) as follows: (a) to compete with AA choosing a mixed strategy $d_{(1,2)}^l = 0,8923$ for a lucrative international route and the rest for $d_{(1,1)}^l = 0.1077$ a domestic route seriously considered by AA, (b) to strongly compete with AA in strategy $d_{(2,1)}^l$ for an international route, and (c) to fully compete with AA in strategies $d_{(3,1)}^l$ and $d_{(4,1)}^l$ for the most lucrative domestic routes. AII, instead, launch a strategy (44) for local and short-haul routes to stay at the forefront of competition deciding (a) to compete with AA in strategy $d_{(1,1)}^l$ for a domestic route, (b) to compete partially with AA choosing a mixed strategy $d_{(2,2)}^l$ = 0.0786 for a lucrative international route and the rest for $d_{(2,1)}^l = 0.9214$ a local route not seriously considered by AA, (c) to absolutely compete with AA in strategy $d_{(3,2)}^l$ choosing seriously only an international route, and (d) to partially compete with AA choosing a mixed strategy $d_{(4,2)}^l = 0.1144$ for a lucrative international route and the rest for $d_{(4,1)}^l = 0.8856$ a local route partially considered by AA.

Remark 2. Price efficiency and effectiveness are critical for an airline's ability to compete and survive. However, this does not mean that every airline should seek to offer the lowest price. Instead, it is important to offer the

appropriate price for routes not considered by the leader airline.



Fig. 2. Convergence of the parameter λ .



Fig. 3. Convergence of the strategies of AA.



Fig. 4. Convergence of the strategies of AI.



Fig. 5. Convergence of the strategies of AII.

7. Conclusion

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The main contribution of this paper was the development of the extraproximal method for computing the Stackelberg/Nash equilibria in a class of ergodic controlled finite Markov chains games. The nonlinear programming problem was represented using an implementation of the Lagrange principle. Another important contribution was the use of the regularizing parameter, which provides strong convexity for the cost-functions and, hence, the correctness of the convergence analysis.

For solving the extraproximal method, we presented a two-step iterated procedure which involved an iterative solution of a quadratic programming problem for the solution of the Stackelberg game in terms of Markov chains. A numerical method was presented for computing the first step of the extraproximal method (parameter λ). The convergence of the suggested procedure to the Stackelberg/Nash equilibrium was also analyzed. It is important to note that we provided all the details needed to implement the extraproximal method in an efficient and numerically stable way for ergodic controlled finite Markov chains games.

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